

Time correlated quantum amplitude damping channel

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Abstract

We analyze the problem of sending classical information through qubit channels where successive uses of the channel are correlated. This work extends the analysis of C. Macchiavello and G. M. Palma to the case of a non-Pauli channel - the amplitude damping channel. Using the channel description outlined in S. Daffer, *et al.*, we derive the correlated amplitude damping channel. We obtain a similar result to C. Macchiavello and G. M. Palma, that is, that under certain conditions on the degree of channel memory, the use of entangled input signals may enhance the information transmission compared to the use of product input signals.

In real physical quantum transmission channels, it is common to have correlated noise acting on consecutive uses. This is in contrast to the class of memoryless channels, in which uncorrelated (independent) noise acts on each use. The problem of the classical capacity of quantum channels with time correlated noise was first considered by C. Macchiavello and G. M. Palma [2]. They analyzed the specific case of sending qubits (quantum states belonging to 2-dimensional Hilbert spaces, each spanned by orthonormal vectors $\{|0\rangle, |1\rangle\}$) with two consecutive uses of a quantum depolarizing channel with partial memory. The action of such a channel on an input state described by the density operator π , is defined by the completely positive, trace-preserving map to an output state, another density operator ρ [1, 2]:

$$\pi \longrightarrow \rho = \Phi(\pi) = (1 - \mu) \sum_{i,j=0}^3 A_{ij}^u \pi A_{ij}^{u\dagger} + \mu \sum_{k=0}^3 A_{kk}^c \pi A_{kk}^{c\dagger},$$

where $0 \leq \mu \leq 1$. With probability $(1 - \mu)$, the noise is uncorrelated and completely specified by the Kraus operators

$$\begin{aligned} A_{00}^u &= p_0 I \otimes I, \quad A_{01}^u = \sqrt{p_0 p_1} I \otimes \sigma_x, \quad A_{02}^u = \sqrt{p_0 p_2} I \otimes \sigma_y, \quad A_{03}^u = \sqrt{p_0 p_3} I \otimes \sigma_z, \\ A_{10}^u &= \sqrt{p_0 p_1} \sigma_x \otimes I, \quad A_{11}^u = p_1 \sigma_x \otimes \sigma_x, \quad A_{12}^u = \sqrt{p_1 p_2} \sigma_x \otimes \sigma_y, \quad A_{13}^u = \sqrt{p_1 p_3} \sigma_x \otimes \sigma_z, \\ A_{20}^u &= \sqrt{p_0 p_2} \sigma_y \otimes I, \quad A_{21}^u = \sqrt{p_1 p_2} \sigma_y \otimes \sigma_x, \quad A_{22}^u = p_2 \sigma_y \otimes \sigma_y, \quad A_{23}^u = \sqrt{p_2 p_3} \sigma_y \otimes \sigma_z, \\ A_{30}^u &= \sqrt{p_0 p_3} \sigma_z \otimes I, \quad A_{31}^u = \sqrt{p_1 p_3} \sigma_z \otimes \sigma_x, \quad A_{32}^u = \sqrt{p_2 p_3} \sigma_z \otimes \sigma_y, \quad A_{33}^u = p_3 \sigma_z \otimes \sigma_z, \end{aligned}$$

while with probability μ , the noise is correlated and specified by

$$A_{00}^c = \sqrt{p_0} I \otimes I, \quad A_{11}^c = \sqrt{p_1} \sigma_x \otimes \sigma_x, \quad A_{22}^c = \sqrt{p_2} \sigma_y \otimes \sigma_y, \quad A_{33}^c = \sqrt{p_3} \sigma_z \otimes \sigma_z.$$

Here, $0 \leq p \leq 1$, $p_0 = (1 - p)$, $p_1 = p_2 = p_3 = \frac{1}{3}p$, and

$$I = \begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix}, \quad \sigma_x = \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix}, \quad \sigma_y = \begin{pmatrix} 0 & -i \\ i & 0 \end{pmatrix}, \quad \sigma_z = \begin{pmatrix} 1 & 0 \\ 0 & -1 \end{pmatrix}$$

are the identity and Pauli matrices respectively. The key object in their analysis is the mutual information for two uses of the channel,

$$I_2(\mathcal{E}) = S(\rho) - \sum_i q_i S(\rho_i) \quad (1)$$

where $\mathcal{E} = \{q_i, \pi_i\}$, with $q_i \geq 0$, $\sum_i q_i = 1$, is the input ensemble of states π_i , transmitted with *a priori* probabilities q_i , of two - generally entangled - qubits. In Eq.(1),

$$S(\sigma) = -\text{tr}(\sigma \log_2 \sigma)$$

is the von Neumann entropy, $\rho = \sum_i q_i \rho_i$, and $\rho_i = \Phi(\pi_i)$. For the following choice of equally-weighted ensemble of orthonormal input states,

$$\begin{aligned}\pi &= \frac{1}{4}(|\pi_1\rangle\langle\pi_1| + |\pi_2\rangle\langle\pi_2| + |\pi_3\rangle\langle\pi_3| + |\pi_4\rangle\langle\pi_4|), \\ |\pi_1\rangle &= \cos\theta|00\rangle + \sin\theta|11\rangle, \\ |\pi_2\rangle &= \sin\theta|00\rangle - \cos\theta|11\rangle, \\ |\pi_3\rangle &= \cos\theta|01\rangle + \sin\theta|10\rangle, \\ |\pi_4\rangle &= \sin\theta|01\rangle - \cos\theta|10\rangle,\end{aligned}\tag{2}$$

Eq.(1) reduces to

$$I_2 = 2 + \sum_{i=1}^4 e_i \log_2 e_i$$

with

$$\begin{aligned}\eta &= 1 - \frac{4}{3}p, \\ e_1 = e_2 &= \frac{1}{4}(1 - \eta^2)(1 - \mu), \\ e_{3,4} &= \frac{1}{4} \left[(1 + \mu) + \eta^2(1 - \mu) \pm 2\sqrt{\eta^2 \cos^2 2\theta + [\mu + \eta^2(1 - \mu)]^2 \sin^2 2\theta} \right].\end{aligned}$$

They showed that there exists a threshold value

$$\mu_t = \frac{\eta}{1 + \eta}$$

for $0 < \eta < 1$, such that when $\mu > \mu_t$, I_2 is maximal for $\theta = \frac{\pi}{4}$ (i.e., the maximally entangled Bell states), while when $\mu < \mu_t$, I_2 is maximal for $\theta = 0$ (i.e., the completely unentangled product states). Furthermore, at $\mu = \mu_t$, any set of states of the form Eq.(2) leads to the same I_2 .

We note that, in the same manner, a quantum dephasing channel with uncorrelated noise can be defined as one specified by the following Kraus operators

$$\begin{aligned}D_{00}^u &= (1 - p)I \otimes I, \quad D_{01}^u = \sqrt{p(1 - p)}I \otimes \sigma_z \\ D_{10}^u &= \sqrt{p(1 - p)}\sigma_z \otimes I, \quad D_{11}^u = p\sigma_z \otimes \sigma_z\end{aligned}\tag{3}$$

and one with correlated noise by

$$D_{00}^c = \sqrt{1 - p}I \otimes I, \quad D_{11}^c = \sqrt{p}\sigma_z \otimes \sigma_z.\tag{4}$$

The same prescription can be applied to a quantum amplitude damping channel with uncorrelated noise as one specified by

$$E_{00}^u = E_0 \otimes E_0, \quad E_{01}^u = E_0 \otimes E_1, \quad E_{10}^u = E_1 \otimes E_0, \quad E_{11}^u = E_1 \otimes E_1, \quad (5)$$

where, with $0 \leq \chi \leq \frac{\pi}{2}$,

$$E_0 = \begin{pmatrix} \cos \chi & 0 \\ 0 & 1 \end{pmatrix}, \quad E_1 = \begin{pmatrix} 0 & 0 \\ \sin \chi & 0 \end{pmatrix} \quad (6)$$

are the Kraus operators for an amplitude damping channel. Here, $|0\rangle$ and $|1\rangle$ denote the excited and ground states respectively. However, it is not *a priori* clear how the Kraus operators for a quantum amplitude damping channel with correlated noise could be constructed in a similar manner, if it is at all possible.

Recently, S. Daffer, *et al.* [3] used a special basis of left, $\{L_i\}$, and right, $\{R_i\}$, damping eigenoperators for a Lindblad superoperator,

$$\Phi(\pi) = \exp(t\mathcal{L})\pi,$$

where t is time, to calculate explicitly the image of a completely positive, trace-preserving map for a wide class of Markov quantum channels:

$$\pi \longrightarrow \rho = \Phi(\pi) = \sum_i \text{tr}(L_i \pi) \exp(\lambda_i t) R_i. \quad (7)$$

For a finite N -dimensional Hilbert space,

$$\mathcal{L}\pi = -\frac{1}{2} \sum_{i,j=1}^{N^2-1} c_{ij} (F_j^\dagger F_i \pi + \pi F_j^\dagger F_i - 2F_i \pi F_j^\dagger),$$

with the system operators F_i satisfying

$$\text{tr}(F_i) = 0, \quad \text{tr}(F_i^\dagger F_j) = \delta_{ij}.$$

The complex c_{ij} form a positive matrix. The right eigenoperators R_i satisfy the eigenvalue equation

$$\mathcal{L}R_i = \lambda_i R_i, \quad (8)$$

and the following duality relation

$$\text{tr}(L_i R_j) = \delta_{ij} \quad (9)$$

with the left eigenoperators L_i . The amplitude damping and dephasing channels are examples of quantum Markov channels. The Lindblad equation [3]

$$\mathcal{L}\pi = -\frac{1}{2}\alpha(\sigma^\dagger\sigma\pi + \pi\sigma^\dagger\sigma - 2\sigma\pi\sigma^\dagger), \quad (10)$$

where α is a parameter analogous to the Einstein coefficient of spontaneous emission, and

$$\sigma^\dagger \equiv \frac{1}{2}(\sigma_x + i\sigma_y), \quad \sigma \equiv \frac{1}{2}(\sigma_x - i\sigma_y)$$

are the creation and annihilation operators respectively, yields the amplitude damping channel, Eq.(6). And, the dephasing channel can be derived from [3]

$$\mathcal{L}\pi = -\frac{1}{2}\Gamma(\pi - \sigma_z\pi\sigma_z), \quad (11)$$

where Γ is another parameter. In this paper, we derive Eq.(4) from the Lindblad equation Eq.(12), that is, it gives the quantum dephasing channel with correlated noise. In a similar fashion, we solve Eq.(19) and interpret the resulting completely positive, trace-preserving map as one which describes a quantum amplitude damping channel with correlated noise. We then analyze, as in Ref.[2], the action of a quantum amplitude damping channel with partial memory on Eq.(2). Our results are in agreement with those of Ref.[2]. That is, the transmission of classical information can be enhanced by employing maximally entangled states as carriers of information rather than product states.

We begin by solving the following Lindblad equation

$$\mathcal{L}\pi = -\frac{1}{2}\Gamma[\pi - (\sigma_z \otimes \sigma_z)\pi(\sigma_z \otimes \sigma_z)]. \quad (12)$$

Eq.(12) is an obvious extension of Eq.(11) with σ_z replaced by $(\sigma_z \otimes \sigma_z)$. The rationale is so that the phase-flip actions of the channel would then be correlated. The method of solution involves first determining the right eigenoperators R_i , which solves Eq.(8):

$$R_{00} = \frac{1}{\sqrt{2}} \begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 \end{pmatrix}, \quad \lambda_{00} = 0; \quad R_{33} = \frac{1}{\sqrt{2}} \begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & -1 \end{pmatrix}, \quad \lambda_{33} = 0; \quad (13)$$

$$R_{01}^\pm = \frac{1}{\sqrt{2}} \begin{pmatrix} 0 & \pm 1 & 0 & 0 \\ 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \end{pmatrix}, \quad \lambda_{01}^\pm = -\Gamma; \quad R_{02}^\pm = \frac{1}{\sqrt{2}} \begin{pmatrix} 0 & 0 & \pm 1 & 0 \\ 0 & 0 & 0 & 0 \\ 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \end{pmatrix}, \quad \lambda_{02}^\pm = -\Gamma; \quad (14)$$

$$R_{03}^{\pm} = \frac{1}{\sqrt{2}} \begin{pmatrix} 0 & 0 & 0 & \pm 1 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 1 & 0 & 0 & 0 \end{pmatrix}, \quad \lambda_{03}^{\pm} = 0; \quad R_{11} = \begin{pmatrix} 0 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \end{pmatrix}, \quad \lambda_{11} = 0; \quad (15)$$

$$R_{12}^{\pm} = \frac{1}{\sqrt{2}} \begin{pmatrix} 0 & 0 & 0 & 0 \\ 0 & 0 & \pm 1 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 \end{pmatrix}, \quad \lambda_{12}^{\pm} = 0; \quad R_{13}^{\pm} = \frac{1}{\sqrt{2}} \begin{pmatrix} 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & \pm 1 \\ 0 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \end{pmatrix}, \quad \lambda_{13}^{\pm} = -\Gamma; \quad (16)$$

$$R_{22} = \begin{pmatrix} 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 \end{pmatrix}, \quad \lambda_{22} = 0; \quad R_{23}^{\pm} = \frac{1}{\sqrt{2}} \begin{pmatrix} 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & \pm 1 \\ 0 & 0 & 1 & 0 \end{pmatrix}, \quad \lambda_{23}^{\pm} = -\Gamma. \quad (17)$$

Notice that the index i here is a doublet. Next, the left eigenoperators are determined by imposing Eq.(9). Finally, the image of the completely positive, trace-preserving map can be obtained via Eq.(7). In this case, we have

$$\pi \longrightarrow \rho = \sum_i \text{tr}(L_i \pi) \exp(\lambda_i t) R_i = \sum_{j=0}^1 D_{jj}^c \pi D_{jj}^{c\dagger},$$

where D_{jj}^c are given by Eq.(4), with

$$p \equiv \frac{1}{2}[1 - \exp(-\Gamma t)]. \quad (18)$$

Therefore, Eq.(12) does indeed yield a dephasing channel with correlated noise.

Next, we solve the Lindblad equation

$$\mathcal{L}\pi = -\frac{1}{2}\alpha[(\sigma^\dagger \otimes \sigma^\dagger)(\sigma \otimes \sigma)\pi + \pi(\sigma^\dagger \otimes \sigma^\dagger)(\sigma \otimes \sigma) - 2(\sigma \otimes \sigma)\pi(\sigma^\dagger \otimes \sigma^\dagger)]. \quad (19)$$

This follows from the same rationale behind the construction of Eq.(12). By replacing σ in Eq.(10) with $(\sigma \otimes \sigma)$, we expect the actions of the resulting channel to be correlated. We call it the quantum amplitude damping channel with correlated noise. The right eigenoperators R_i , which solves Eq.(8) are

$$R_{00} = \frac{1}{\sqrt{2}} \begin{pmatrix} 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 2 \end{pmatrix}, \quad \lambda_{00} = 0; \quad R_{33} = \frac{1}{\sqrt{2}} \begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & -1 \end{pmatrix}, \quad \lambda_{33} = -\alpha \quad (20)$$

and those in Eq.(14) to Eq.(17), but with the following respective eigenvalues,

$$\begin{aligned}\lambda_{01}^{\pm} &= \lambda_{02}^{\pm} = \lambda_{03}^{\pm} = -\frac{1}{2}\alpha, \\ \lambda_{11} &= \lambda_{12}^{\pm} = \lambda_{13}^{\pm} = \lambda_{22} = \lambda_{23}^{\pm} = 0.\end{aligned}\tag{21}$$

The left eigenoperators are determined as above, and Eq.(7) becomes

$$\pi \longrightarrow \rho = \sum_i \text{tr}(L_i \pi) \exp(\lambda_i t) R_i = \sum_{j=0}^1 E_{jj}^c \pi E_{jj}^{c\dagger},$$

where

$$E_{00}^c = \begin{pmatrix} \cos \chi & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \end{pmatrix}, \quad E_{11}^c = \begin{pmatrix} 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ \sin \chi & 0 & 0 & 0 \end{pmatrix}, \tag{22}$$

with

$$\cos \chi \equiv \exp\left(-\frac{1}{2}\alpha t\right), \quad \sin \chi \equiv \sqrt{1 - \exp(-\alpha t)}.\tag{23}$$

We note that, in contrast to D_{00}^c in Eq.(4), E_{00}^c cannot be written as a tensor product of two two-by-two matrices. This gives rise to the typical “spooky” action of the channel: $|01\rangle$, $|10\rangle$, and any linear combination of them, and $|11\rangle$ will go through the channel undisturbed, but not $|00\rangle$.

It is interesting to note that Eq.(3) can also be derived by solving the following Lindblad equation,

$$\mathcal{L}\pi = -\frac{1}{2}\Gamma[\pi - (I \otimes \sigma_z)\pi(I \otimes \sigma_z)] - \frac{1}{2}\Gamma[\pi - (\sigma_z \otimes I)\pi(\sigma_z \otimes I)].$$

However, analogous approach for the amplitude damping channel with uncorrelated noise does not work. This is because the amplitude damping channel is by definition non-unital. This is not surprising in view of the fact that although all Lindblad superoperators have a Kraus decomposition, the converse is not true in general.

Now, we carry out the same analysis as in Ref.[2] by considering

$$\pi \longrightarrow \rho = \Phi(\pi) = (1 - \mu) \sum_{i,j=0}^1 E_{ij}^u \pi E_{ij}^{u\dagger} + \mu \sum_{i=0}^1 E_{ii}^c \pi E_{ii}^{c\dagger} \tag{24}$$

and using Eq.(2). Eq.(1) then becomes

$$I_2 = -\sum_{i=1}^4 t_i \log_2 t_i + \frac{1}{4} \sum_{j=1}^4 u_j \log_2 u_j + \frac{1}{4} \sum_{k=1}^4 v_k \log_2 v_k + \frac{1}{2} \sum_{l=1}^2 w_l \log_2 w_l$$

with

$$\begin{aligned}
\Theta(\mu, \chi) &\equiv \frac{1}{2} \left[(3 + \mu) + (1 - \mu) \left(\cos 4\chi - 32\mu \cos^2 \chi \sin^4 \frac{\chi}{2} \right) \right], \\
t_{1,2} &= \frac{1}{4} (1 \pm \sin^2 \chi) [(1 \pm \sin^2 \chi) \mp \mu \sin^2 \chi], \\
t_{3,4} &= \frac{1}{4} [1 - (1 - \mu) \sin^4 \chi], \\
u_{1,2} &= (1 - \mu) \cos^2 \theta \cos^2 \chi \sin^2 \chi, \\
u_{3,4} &= \frac{1}{2} - (1 - \mu) \cos^2 \theta \cos^2 \chi \sin^2 \chi \pm \frac{1}{2} \sqrt{\cos^4 \theta \cos^2 2\chi + \sin^4 \theta + \cos^2 \theta \sin^2 \theta \Theta(\mu, \chi)}, \\
v_{1,2} &= (1 - \mu) \sin^2 \theta \cos^2 \chi \sin^2 \chi, \\
v_{3,4} &= \frac{1}{2} - (1 - \mu) \sin^2 \theta \cos^2 \chi \sin^2 \chi \pm \frac{1}{2} \sqrt{\sin^4 \theta \cos^2 2\chi + \cos^4 \theta + \cos^2 \theta \sin^2 \theta \Theta(\mu, \chi)}, \\
w_1 &= \mu + (1 - \mu) \cos^2 \chi, \quad w_2 = (1 - \mu) \sin^2 \chi.
\end{aligned} \tag{25}$$

Numerical calculation of I_2 for $\theta = \frac{\pi}{4}$ (i.e., the maximally entangled Bell states), and $\theta = 0$ (i.e., the completely unentangled product states), with $0 \leq \chi \leq \frac{\pi}{2}$ and $0 \leq \mu \leq 1$ allows us to compare the information carrying capacity of both forms of input state. Constructing graphs similar to that in Ref.[2] shows for each χ there is a threshold μ_t such that for $\mu > \mu_t$, the performance of the Bell states for classical information transmission is better than that of the product states. While, for $\mu < \mu_t$, better performance is achieved by using the product states instead. For instance, when $\chi = \frac{\pi}{5}$, we have $\mu_t \in (0.5, 0.6)$. Furthermore, for the product states

$$I_2(\theta = 0, \mu = 1, \chi) \geq I_2(\theta = 0, \mu = 0, \chi). \tag{26}$$

This is a direct consequence of the following inequality:

$$a \log a + x \log x \leq (a + x) \log(a + x), \quad \forall a, x \geq 0.$$

What this shows is that in the case of quantum amplitude damping channel with perfect memory, it is possible to obtain enhanced information carrying performance even if we are using product input states. This is the same conclusion reached in Ref.[2] for the case of a depolarizing channel.

In conclusion, we have extended the problem of time-correlated noise (or “channels with memory”) as considered in Ref.[2] to the case of the amplitude damping channel. In the case of sending two qubits by successive uses of an amplitude damping channel with partial memory, we establish numerically that by using maximally entangled states rather than product

states as information carriers, we can enhance the transmission of classical information over the quantum channel.

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